On the Expressive Power of Deep Learning: A Tensor Analysis

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The Expressive Power of Deep Learning

- Expressive power of depth the driving force behind Deep Learning
- **Depth efficiency**: when a polynomially sized deep network realizes a function that requires shallow networks to have super-polynomial size
- Prior works on depth efficiency:
 - Show its existence, without discussing how frequent it is
 - Do not apply to convolutional networks (locality+sharing+pooling), the most successful deep learning architecture to date

Convolutional Arithmetic Circuits



Convolutional networks:

- locality
- sharing (optional)
- product pooling

Computation in log-space leads to SimNets – new deep learning architecture showing promising empirical performance ¹

Cohen, Sharir, Shashua (HUJI) On the Expressive Power of Deep Learning

¹Deep SimNets, CVPR'16

Convolutional Arithmetic Circuits (cont')



Function realized by output *y*:

$$h_{\mathcal{Y}}(\mathbf{x}_{1},\ldots,\mathbf{x}_{N}) = \sum_{d_{1}\ldots d_{N}=1}^{M} \mathcal{A}_{d_{1},\ldots,d_{N}}^{\mathcal{Y}} \prod_{i=1}^{N} f_{\theta_{d_{i}}}(\mathbf{x}_{i})$$

- **x**₁...**x**_N input patches
- $f_{\theta_1} \dots f_{\theta_M}$ representation layer functions

• \mathcal{A}^{y} - coefficient tensor (M^{N} entries, polynomials in weights $\mathbf{a}^{l,j,\gamma}$)

Shallow Network \leftrightarrow CP Decomposition

Shallow network (single hidden layer, global pooling):



Coefficient tensor $\mathcal{A}^{\mathcal{Y}}$ given by classic **CP decomposition**:

$$\mathcal{A}^{y} = \sum_{\gamma=1}^{r_{0}} a_{\gamma}^{1,1,y} \cdot \underbrace{\mathbf{a}^{0,1,\gamma} \otimes \mathbf{a}^{0,2,\gamma} \otimes \cdots \otimes \mathbf{a}^{0,N,\gamma}}_{\text{rank-1 tensor}} (rank(\mathcal{A}^{y}) \leq r_{0})$$

Deep Network \leftrightarrow Hierarchical Tucker Decomposition

Deep network ($L = \log_2 N$ hidden layers, size-2 pooling windows):



Coefficient tensor $\mathcal{A}^{\mathcal{Y}}$ given by **Hierarchical Tucker decomposition**:

$$\begin{split} \phi^{1,j,\gamma} &= \sum_{\alpha=1}^{\prime_0} a_{\alpha}^{1,j,\gamma} \cdot \mathbf{a}^{0,2j-1,\alpha} \otimes \mathbf{a}^{0,2j,\alpha} \\ & \cdots \\ \phi^{l,j,\gamma} &= \sum_{\alpha=1}^{\prime_{l-1}} a_{\alpha}^{l,j,\gamma} \cdot \phi^{l-1,2j-1,\alpha} \otimes \phi^{l-1,2j,\alpha} \\ & \cdots \\ \mathcal{A}^{y} &= \sum_{\alpha=1}^{\prime_{l-1}} a_{\alpha}^{l,1,y} \cdot \phi^{l-1,1,\alpha} \otimes \phi^{l-1,2,\alpha} \end{split}$$

Theorem of Network Capacity

Theorem

The rank of tensor \mathcal{A}^{y} given by Hierarchical Tucker decomposition is at least min $\{r_0, M\}^{N/2}$ almost everywhere w.r.t. decomposition parameters.

Since rank of $\mathcal{A}^{\mathcal{Y}}$ generated by CP decomposition is no more than the number of terms (# of hidden channels in shallow network):

Corollary

Randomizing linear weights of deep network by a continuous distribution gives functions that **with probability one**, cannot be approximated by shallow network with less than $\min\{r_0, M\}^{N/2}$ hidden channels.

Depth efficiency holds almost always!

Theorem of Network Capacity – Proof Sketch

- $\llbracket \mathcal{A} \rrbracket$ arrangement of tensor \mathcal{A} as matrix (*matricization*)
- \odot Kronecker product for matrices. Holds: $rank(A \odot B) = rank(A) \cdot rank(B)$
- Relation between tensor and Kronecker products: $\llbracket \mathcal{A} \otimes \mathcal{B} \rrbracket = \llbracket \mathcal{A} \rrbracket \odot \llbracket \mathcal{B} \rrbracket$
- Implies: $\mathcal{A} = \sum_{z=1}^{Z} \lambda_z \mathbf{v}_1^{(z)} \otimes \cdots \otimes \mathbf{v}_{2^L}^{(z)} \Longrightarrow rank \llbracket \mathcal{A} \rrbracket \leq Z$
- By induction over l = 1...L, almost everywhere w.r.t. $\{\mathbf{a}^{l,j,\gamma}\}_{l,j,\gamma}$: $\forall j \in [N/2'], \gamma \in [r_l] : rank \llbracket \phi^{l,j,\gamma} \rrbracket \ge (\min\{r_0, M\})^{2'/2}$
 - Base: "SVD has maximal rank almost everywhere"
 - <u>Step</u>: $rank[\![\mathcal{A} \otimes \mathcal{B}]\!] = rank([\![\mathcal{A}]\!] \odot [\![\mathcal{B}]\!]) = rank[\![\mathcal{A}]\!] \cdot rank[\![\mathcal{B}]\!]$, and "linear combination preserves rank almost everywhere"

Generalization

Comparison between arbitrary depths shows penalty in resources grows *double exponentially* w.r.t. number of layers cut off.



Through tensor decompositions, we showed that *depth efficiency holds almost always with convolutional arithmetic circuits*

Equivalence between convolutional networks and tensor decompositions has many other applications, for example:

- Expressiveness of convolutional ReLU networks: ¹
 - Average pooling leads to loss of universality
 - Depth efficiency exists but does not hold almost always
- Inductive bias of convolutional arithmetic circuits: ²
 - Deep networks can model strong correlation between input elements, shallow networks can't
 - Pooling geometry of a deep network selects supported correlations

¹Convolutional Rectifier Networks as Generalized Tensor Decompositions, ICML'16 ²Inductive Bias of Deep Convolutional Networks through Pooling Geometry, arXiv

Thank You